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DATA PROCESSING SYSTEMS

2 FIELD OF INVENTION

3 The present invention is directed to data processing
4 systems. More particularly, it is directed to systems,
5 methods and apparatus for introducing security measures to
6 data processing systems.

7 BACKGROUND OF THE INVENTION

8 A data processing system typically comprises a central
9 processing unit (CPU), an input/output (I/O) subsystem, and
10 a memory subsystem, all interconnected by a bus subsystem.
11 The memory subsystem typically comprises random access
12 memory (RAM), read only memory (ROM), and one or more data
13 storage devices such as hard disk drives, optical disk
14 drives, and the like. The I/O subsystem typically comprises:
15 a display; a printer; a keyboard; a pointing device such as
16 a mouse, tracker ball, or the like; and one or more network
17 connections permitting communications between the data
18 processing system and one or more similar systems and/or
19 peripheral devices via a data communications network. The
20 combination of such systems and devices interconnected by
21 such a network may itself form a distributed data processing
22 system. Such distributed systems may be themselves
23 interconnected by additional data communications networks.
24 In the memory subsystem is stored data and computer program
25 code executable by the CPU. The program code includes
26 operating system software and application software. The
27 operating system software, when executed by the CPU,
28 provides a platform on which the application software can be

1 executed. The operating system software has a core or kernel
2 of code in which the basic functions of the operating system
3 software are defined.

4 A problem associated with data processing systems is that of
5 security. In particular, it is becoming increasingly
6 difficult to determine with any degree of certainty that a
7 data processing system actually has the properties it is
8 believed to have. This difficulty arises because data
9 processing systems, and particularly the operating systems
10 therein, are becoming increasingly general purpose,
11 configurable, and reconfigurable in nature. The
12 administrative state of a data processing systems can be
13 varied from one moment to the next based on an
14 administrative action. Specifically, the administrative
15 state of a data processing system is defined by the
16 combination of software and data present in the machine. The
17 software may include binary files, patches, applications,
18 and the like added to and deleted from the system from time
19 to time via one or more administrative actions. An
20 administrative action such as the addition or deletion of
21 software in the system can thus be regarded as a change in
22 the state of the system. Many data processing systems can be
23 placed into a corrupt state by users and/or system
24 administrators with or without proper authorization. This
25 form of corruption is difficult to detect. It would be
26 desirable to make such corruption easier to detect.

27 Many data processing networks employ intrusion detection and
28 diagnosis (IDD) systems. These IDD systems are typically
29 data processing systems resident on the network and
30 dedicated to intrusion detection and diagnosis. It will be
31 appreciated that detection of corruption is important in the
32 field of IDD. Most intruders do not want to be detected.
33 Thus, administration tools employed in IDD systems are among
34 the first to be attacked.

1 Cracker tools allow hackers, crackers, or other attackers to
2 selectively hide files, processes, and network connections
3 in an individual host data processing system. An example of
4 a conventional cracker tool is known as "rootkit". Rootkit
5 replaces Unix system commands used for investigation, such
6 as ls, ps, netstat, and ifconfig, with so-called Trojan
7 horse versions that hide the activities of an attacker.
8 Conventionally, such Trojan horses have been identified by
9 calculating, storing, and comparing databases of
10 cryptographic checksums of system binaries. However, recent
11 versions of "rootkit" include Trojan horse versions of the
12 programs employed to generate and compare the checksums.
13 Attackers have recently begun to employ loadable kernel
14 modules to introduce Trojan horses to data processing
15 systems. A kernel is difficult to inspect when running.
16 Thus, Trojan horse modules therein remain undetected. A
17 typical defense against such Trojan horse modules is to
18 prevent kernel modules from being loaded. However, system
19 functionality is then limited.

20 There is growing interest in releasing computer software in
21 source code form under "open source" licenses such as the
22 "General Public License" and the like. Such releases
23 facilitate creation of Trojan horses, particularly when the
24 software in question is operating system software. The
25 detection of Trojan horses is therefore of increasing
26 importance. IDD systems are an early target for infiltration
27 by Trojan horses. Here, the attacker typically alters the
28 IDD system in such a manner that it appears to continue
29 functioning normally, but in reality hides the attacker's
30 activities.

31 Conventional security schemes for data processing systems
32 include secure logging schemes and forward secrecy of
33 digital signatures.

1 Secure logging schemes are directed to the protection of
2 integrity, secrecy, and authenticity of records of data
3 processing events. The schemes may be employed in
4 maintaining quality standards of event logging. In general,
5 secure logging schemes assume the existence of a secure
6 logging host data processing system and operate in
7 dependence on a combination of message or stream encryption,
8 hash chaining, authentication codes, and one way key
9 evolution.

10 Forward secrecy of digital signatures is directed to
11 limiting damage to compromised signature keys. In operation,
12 forward secrecy of digital signatures provides a series of
13 signing keys:

14 $SK_0 \Rightarrow SK_1 \Rightarrow SK_2 \Rightarrow \dots$

15 So that SK_{n+1} is a derivation of SK_n , and that verification
16 of a signature does not require distribution, traversal, and
17 verification of a chain of keys.

18 **SUMMARY OF THE INVENTION**

19 In accordance with the present invention, there is now
20 provided systems apparatus and methods for detecting an
21 attack on a data processing system. An example method
22 comprising, in the data processing system: providing an
23 initial secret; binding the initial secret to data
24 indicative of an initial state of the system via a
25 cryptographic function; recording state changing
26 administrative actions performed on the system in a log;
27 prior to performing each state changing administrative
28 action, generating a new secret by performing the

1 cryptographic function on a combination of data indicative
2 of the administrative action and the previous secret, and
3 erasing the previous secret; evolving the initial secret
4 based on the log to produce an evolved secret; comparing the
5 evolved secret with the new secret; determining that the
6 system is uncorrupted if the comparison indicates a match
7 between the evolved secret and the new secret; and,
8 determining that the system is corrupted if the comparison
9 indicate a mismatch between the evolved secret and the new
10 secret.

11 From another aspect, there is now also provided a data
12 processing system comprising: a processor; a memory
13 connected to the processor; and, detection logic connected
14 to the processor and the memory, the detection logic, in
15 use: providing an initial secret; binding the initial secret
16 to data indicative of an initial state of the system via a
17 cryptographic function; recording state changing
18 administrative actions performed on the system in a log;
19 prior to performing each state changing administrative
20 action, generating a new secret by performing the
21 cryptographic function on a combination of data indicative
22 of the administrative action and the previous secret, and
23 erasing the previous secret; evolving the initial secret
24 based on the log to produce an evolved secret; comparing the
25 evolved secret with the new secret; determining that the
26 system is uncorrupted if the comparison indicates a match
27 between the evolved secret and the new secret; and,
28 determining that the system is corrupted if the comparison
29 indicate a mismatch between the evolved secret and the new
30 secret.

31 The present invention advantageously also provides methods
32 for cryptographic entangling of state and administration in
33 a data processing system that permits detection of Trojan
34 horses within the system.

1 **BRIEF DESCRIPTION OF THE DRAWINGS**

2 The invention and its embodiments will be more fully
3 appreciated by reference to the following detailed
4 description of advantageous and illustrative embodiments in
5 accordance with the present invention when taken in
6 conjunction with the accompanying drawings, in which:

7 Fig. 1 shows a block diagram of an example of a data
8 processing system embodying the present invention;

9 Fig. 2 is a flow diagram associated with the data processing
10 system;

11 Fig. 3 is another flow diagram associated with the data
12 processing system;

13 Fig. 4 is yet another flow diagram associated with the data
14 processing system;

15 Fig. 5 is a further flow diagram associated with the data
16 processing system; and,

17 Fig. 6 is a block diagram of a data processing system in a
18 layer format.

19 **DETAILED DESCRIPTION OF THE INVENTION**

20 The present invention provides methods, systems and
21 apparatus for detecting an attack on a data processing
22 system. An example method comprising, in the data

1 processing system: providing an initial secret; binding the
2 initial secret to data indicative of an initial state of the
3 system via a cryptographic function; recording state
4 changing administrative actions performed on the system in a
5 log; prior to performing each state changing administrative
6 action, generating a new secret by performing the
7 cryptographic function on a combination of data indicative
8 of the administrative action and the previous secret, and
9 erasing the previous secret; evolving the initial secret
10 based on the log to produce an evolved secret; comparing the
11 evolved secret with the new secret; determining that the
12 system is uncorrupted if the comparison indicates a match
13 between the evolved secret and the new secret; and,
14 determining that the system is corrupted if the comparison
15 indicate a mismatch between the evolved secret and the new
16 secret. The term 'in the system' as used herein, includes
17 elements peripheral to the system but coupled thereto
18 directly or indirectly.

19 The cryptographic function advantageously comprises a
20 one-way hash function. The hash function may comprise an
21 exponentiation function. In an advantageous embodiment of
22 the present invention, the cryptographic function comprises
23 a public/private key pair. The initial secret may be
24 received from a system administrator.

25 Viewing the present invention from another aspect, there is
26 now provided a data processing system comprising: a
27 processor; a memory connected to the processor; and,
28 detection logic connected to the processor and the memory,
29 the detection logic, in use: providing an initial secret;
30 binding the initial secret to data indicative of an initial
31 state of the system via a cryptographic function; recording
32 state changing administrative actions performed on the
33 system in a log; prior to performing each state changing
34 administrative action, generating a new secret by performing

1 the cryptographic function on a combination of data
2 indicative of the administrative action and the previous
3 secret, and erasing the previous secret; evolving the
4 initial secret based on the log to produce an evolved
5 secret; comparing the evolved secret with the new secret;
6 determining that the system is uncorrupted if the comparison
7 indicates a match between the evolved secret and the new
8 secret; and, determining that the system is corrupted if the
9 comparison indicate a mismatch between the evolved secret
10 and the new secret.

11 The present invention also extends to a computer program
12 element comprising computer program code means which, when
13 loaded in a processor of a computer system, configures the
14 processor to perform the aforementioned attack detection
15 method.

16 The present invention advantageously provides a method for
17 cryptographic entangling of state and administration in a
18 data processing system that permits detection of Trojan
19 horses within the system.

20 In an advantageous embodiment of the present invention, an
21 initial secret is evolved in a data processing system in an
22 irreversible manner. The evolution progresses based on
23 administrative actions in the system. As the evolution
24 progresses, previous secrets are overwritten. Proof of
25 knowledge of the evolved secret thus equates to a
26 cryptographic verification of the history of administrative
27 action in the system. Thus, if the initial state of the
28 system is known, the current state can be determined. If the
29 initial state is known to have been an uncorrupted state,
30 then an absence of subsequent corruption can be proved. It
31 can then be determined that the system remains uncorrupted.

1 In a particularly advantageous embodiment of the present
2 invention, the aforementioned cryptographic entangling has
3 initialization, update, and proof stages. In the
4 initialization stage, the system generates an initial secret
5 and releases binding data that binds to the secret. The
6 binding data may take different forms depending on
7 application.

8 In the update stage, the system updates the secret each time
9 there is an administrative action that could lead to system
10 corruption such as infiltration of system level Trojan
11 horses. Examples of such corruption actions include:
12 updating of system executable code; updating of system
13 libraries; installation of kernel modules; reading of files
14 such as those used to store system states during rebooting
15 operations; alteration of configuration files; alteration of
16 system run-level codes; and, writing to or reading from
17 peripheral devices. Each update proceeds in the following
18 manner. First, a new secret is computed by applying a one
19 way function to the combination of the previous secret, a
20 description of the action, and associated context
21 information. Then, the previous secret and any information
22 from which it might be derived is completely overwritten or
23 otherwise erased. A description of the action is then
24 logged. Finally the action is performed.

25 In the proof stage, the system offers a proof that its
26 current secret corresponds to the initial secret as it has
27 evolved according to the record of logged actions.

28 Referring first to Figure 1, a data processing system
29 comprises a CPU 10, an I/O subsystem 20, and a memory
30 subsystem 40, all interconnected by a bus subsystem 30. The
31 memory subsystem 40 may comprise random access memory (RAM),
32 read only memory (ROM), and one or more data storage devices
33 such as hard disk drives, optical disk drives, and the like.

1 The I/O subsystem 20 may comprises: a display; a printer; a
2 keyboard; a pointing device such as a mouse, tracker ball,
3 or the like; and one or more network connections permitting
4 communications between the data processing system and one or
5 more similar systems and/or peripheral devices via a data
6 communications network. The combination of such systems and
7 devices interconnected by such a network may itself form a
8 distributed data processing system. Such distributed systems
9 may be themselves interconnected by additional data
10 communications networks.

11 In the memory subsystem 40 is stored data 60 and computer
12 program code 50 executable by the CPU 10. The program code
13 50 includes operating system software 90 and application
14 software 80. The operating system software 90, when executed
15 by the CPU 10, provides a platform on which the application
16 software 80 can be executed. The operating system software
17 90 has a core or kernel 100 of code in which the basic
18 functions of the operating system software 90 are defined.

19 In an advantageous embodiment of the present invention, the
20 program code 50 stored in the program code 50 stored in the
21 memory subsystem 40 also includes detector logic comprising
22 a cryptographic entangling facility (CEF) 70. The CEF 70 is
23 associated with a log file 110 also stored in the memory
24 subsystem 40. The CEF 70 may be separate from the operating
25 system software 90 as shown or embedded in the operating
26 system software 90. In a particularly advantageous
27 embodiment of the present invention, the CEF 70 is embedded
28 in the kernel 100 of the operating system software 90.

29 In operation, the CEF 70 entangles the state and
30 administration of the data processing system in a
31 cryptographic manner. The entangling permits detection of
32 Trojan horses and the like within the system.

1 As indicated earlier, the administration of the data
2 processing system includes administrative actions such as
3 the addition, deletion, or other reconfiguration of the
4 program code 50 and data 60 recorded in the memory subsystem
5 40. Each administrative action effectively changes the data
6 processing system from one state to another.

7 The cryptographic entangling is developed based on an
8 initial secret. The secret is evolved by the CEF 70 in an
9 irreversible manner. The evolution of the secret progresses
10 in steps. Each step corresponds to and is triggered by a
11 change in state of the data processing system. Whenever an
12 administrative action is to be performed in the data
13 processing system, producing a change of state in the data
14 processing system, there is a corresponding and preceding
15 step in the evolution of the secret. At each step, the CEF
16 70 overwrites or otherwise deletes the previous secret from
17 the memory subsystem 40. The CEF 70 produces a new secret
18 based on the previous secret together with data indicative
19 of the corresponding administrative action. The new secret
20 is recorded by the CEF 70 in place of the previous secret.
21 The administrative action is then performed. Proof of
22 knowledge of the evolved secret thus equates to a
23 cryptographic verification of the history of administrative
24 action in the data processing system. Thus, if the initial
25 state of the system is known, the current state can be
26 determined. If the initial state is known to have been an
27 uncorrupted state, then an absence of subsequent corruption
28 can be proved. It can then be determined that the system
29 remains uncorrupted.

30 Referring now to Figure 2, in advantageous embodiments of
31 the present invention, the cryptographic entangling has an
32 initialization stage 200, an update stage 210, and a proof
33 stage 220.

1 Referring to Figure 3, in the initialization stage 200, the
2 CEF 70, at step 300 generates an initial secret. The initial
3 secret is supplied a system administrator via a secure
4 communication channel. Alternatively, in other embodiments
5 of the present invention, the initial secret may be entered
6 to the CEF 70 by the system administrator. At step 310, the
7 CEF 70 releases binding data. At step 320, the CEF 70 binds
8 the binding data to the secret. The binding data may take
9 different forms depending on the data processing system, the
10 or each application of the data processing system, and trust
11 mechanisms associated with communication of the secret.

12 Referring to Figure 4, in the update stage 210, the CEF 70
13 updates the secret each time there is an administrative
14 action. Specifically, the update is triggered by and
15 performed is advance of administrative action. As indicated
16 earlier, examples of such actions include: updating of
17 system executable code; updating of system libraries;
18 installation of kernel modules; reading of files such as
19 those used to store system states during rebooting
20 operations; alteration of configuration files; alteration of
21 system run-level codes; and, writing to or reading from
22 peripheral devices. Each update proceeds in the following
23 manner. First, at block 400, the CEF 70 computes a new
24 secret. The new secret is computed by the CEF 70 applying a
25 one way function to the combination of the previous secret
26 and data indicative of the administrative action. At block
27 410, the CEF 70 erases the previous secret, together with
28 any information from which it might be derived. At block
29 420, the CEF 70 records data indicative of the
30 administrative action in the log file 110. At block 430, the
31 CEF 70 permits execution of the administrative action.

32 With reference to Figure 5, in the proof stage 220, the CEF
33 70 offers a proof that its current secret corresponds to the
34 initial secret as it has evolved according to the record of

1 actions contained in the log file 110. Specifically, at
2 block 500, the CEF 70 retrieves the initial secret. The
3 initial secret may be retrieved, for example, via a request
4 for entry of the initial secret by a system administrator.
5 At block 510, the CEF 70 retrieves the record of
6 administrative actions from the log file 110 stored in the
7 memory subsystem 40. At block 520, the CEF 70 evolves the
8 initial secret based on the record of administrative actions
9 retrieved from the log file 110. At block 530, the CEF 70
10 compares the secret evolved in block 520 with its current
11 secret. If the secrets match, then, at block 550, the CEF 70
12 reports that the data processing system is still in an
13 uncorrupted state. If however the secrets do not match,
14 then, at block 540, the CEF 70 reports that the data
15 processing system is in a potentially corrupted or otherwise
16 compromised state.

17 The operation of the CEF 70 in a particularly advantageous
18 embodiment of the present invention will now be described.
19 In the following explanation, an ordered collection of
20 administrative actions in the data processing system is
21 denoted by $\{M_i\}$, where i is an index. The index may for
22 example be time. However, such it will appreciated that
23 intervals between successive administrative actions may be
24 irregular. The secret known by the CEF 70 at i is S_i . $P(S_i)$
25 is data that provably binds the data processing system to
26 the secret S_i . During the initialization stage 200, the CEF
27 70 chooses secret S_0 and releases data $P(S_0)$. As indicated
28 earlier, the form of release depends on application. In the
29 update stage 210, the secret S_i is updated according to the
30 mapping $(S_n, M_n) \Rightarrow S_{n+1}$. The mapping $(S_n, M_n) \Rightarrow S_{n+1}$ is a one way
31 collision resistant cryptographic hash function such as MD-5
32 or SHA-1. Collision resistance in this context implies that
33 the pair S_n and S_{n+1} together provide a binding to

1 administrative action M_n . The evolution can thus be
2 represented by:

3 $S_0 \xrightarrow{M_0} S_1 \xrightarrow{M_1} S_2 \dots \xrightarrow{M_{n-1}} S_n$
4 $\downarrow \dots \downarrow \dots \downarrow \dots \downarrow$
5 $P(S_0) \rightarrow^{M_0} P(S_1) \rightarrow^{M_1} P(S_2) \dots \rightarrow^{M_{n-1}} P(S_n)$

6 Where the mapping $P(S_n) \rightarrow P(S_{n+1})$ is an indication that an
7 administrator can verify that data $P(S_{n+1})$ is derived from
8 data $P(S_n)$ according to administrative action M_n .

9 In an advantageous embodiment of the present invention, the
10 initial secret is shared between the data processing system
11 and the administrators of the data processing system, as
12 indicated earlier. The computational overhead associated
13 with this arrangement is advantageously low. However, this
14 arrangement also assumes that all the administrators are
15 benign. At the initialization stage 200, the initial secret
16 S_0 is generated by the CEF 70 and provided via a secure
17 communications channel to the administrators. In the update
18 stage 210, the CEF 70 employs a collision resistant
19 cryptographic hash function \hbar to update the secret, as
20 follows:

21 $S_{n+1} = \hbar("Update", S_n, M_n)$

22 "Update" is a text word included in each update to avoid
23 spoofing attacks.

24 In the proof stage 220, an administrator supplies the index
25 t to the CEF 70. In response, the CEF 70 returns $\hbar(S_n, t)$. The
26 initial secret S_0 is known to administrator. In addition,
27 the history of administrative actions $\{M_n\}$ is recorded in
28 the log file 110. S_n can thus be computed. Because \hbar is

1 collision free, it is possible to conclude that the data
2 processing system knows S_n .

3 In another advantageous embodiment of the present invention,
4 public key encryption is employed. This has an advantage in
5 that the administrators of the data processing system need
6 not trust one another. In the initialization stage 200, the
7 CEF 70 generates a public/private key pair and publishes the
8 public key.

9 In the update stage 210, CEF 70 generates, for each
10 administrative action, a new public/private key pair. The
11 CEF 70 then signs the combination of the new key and data
12 indicative of the administrative action M_n with the
13 previous key. The CEF 70 records both the signature and the
14 signed data in the log file. The CEF 70 then erases the old
15 private key. Only then does the CEF 70 permit the data
16 processing system to perform the administrative action.

17 In the proof stage 220, the administrator supplies index t .
18 In response, the CEF 70 then signs the combination of t with
19 the current public key. The history of administrative
20 actions $\{M_n\}$ is known. Therefore, the chain of signatures
21 can be verified. Therefore, the current public key can be
22 computed. The ability of the CEF 70 to sign a user
23 influenced message demonstrates that the CEF 70 knows the
24 private key. Thus, the state of the data processing system
25 can be confirmed.

26 Another advantageous embodiment of the present invention
27 involves large prime numbers. In the initialization stage
28 200, the CEF 70 chooses two large prime numbers p_1, p_2 . The
29 CEF 70 also chooses an element g of high order in the group
30 $G = (Z/p_1p_2Z)^*$. The CEF 70 picks x where random $x \in [1, p_1p_2]$. In
31 addition, the CEF 70 sets $S_0 = g^x$. The CEF 70 further picks y

1 where random $y \in [1, p_1 p_2]$. The CEF 70 publishes g^y , g , y , and
2 the product $p_1 p_2$. The CEF 70 thus binds the data processing
3 system to g^x .

4 In the update stage 210, the CEF 70 generates new secrets
5 based on $S_{n+1} = (S_n)^{\hbar(n, M_n)} \cdot g$. This operation is effectively one
6 way as it is difficult to take roots in the group G .
7 Multiplication by g is added to avoid degenerately small
8 subgroups.

9 In the proof stage 220, the administrator supplies t . In
10 response, the CEF 70 returns $(S_n)^2 \hbar(t)$. The history of
11 administrative actions $\{M_n\}$ is known from the log file 110.
12 Modular exponentiation inside the group G can be performed.
13 The CEF 70 can compute $(S_n)^y \hbar(t)$ by mimicking the updates and
14 computations that have been performed using g^y rather than
15 the system value of g^x . The value returned raised to the
16 power y can be verified.

17 Yet another particularly advantageous embodiment of the
18 present invention involves Sophie-Germain prime numbers. In
19 the initialization stage 200, the CEF 70 chooses two
20 Sophie-Germain prime numbers p_1, p_2 . The CEF 70 also chooses
21 an element g in the group $G = (Z/p_1 p_2 Z)^x$ with a large prime
22 order greater than the maximum of the hash function $\hbar(\cdot)$. In
23 addition, the CEF 70 chooses random $x \in [1, p_1 p_2]$. The CEF 70
24 also chooses random $y \in [1, p_1 p_2]$. The CEF 70 sets $S_0 = g^x$. In
25 addition, the CEF 70 publishes g^y , g , y , and the product $p_1 p_2$
26 . In the update stage 210, the CEF 70 generates new secrets
27 based on $S_{n+1} = (S_n)^{\hbar(n, M_n)}$. As indicated earlier, it is difficult
28 to take roots in the group G . Thus, $S_{n+1} = (S_n)^{\hbar(n, M_n)}$ is
29 effectively one way. Because $\hbar(n, M_n) < \text{order}(g)$, no entropy is
30 lost. Degenerate subgroups are thus avoided. The proof stage

1 220 here is similar to that described in the preceding
2 example.

3 With reference now to Figure 6, a data processing system as
4 herein before described with reference to Figure 1 can be
5 represented as a stack of data processing layers 600-630.
6 Each layer in the stack forms a foundation of processing for
7 the next layer in ascending order. At the base level, there
8 is the hardware layer 630 comprising the CPU 10, and logic
9 circuitry of the I/O subsystem 20, memory subsystem 40, and
10 bus subsystem 30. Beyond the hardware layer is the kernel
11 layer 620, comprising the kernel 100. Beyond the kernel
12 layer 620 is the operating system layer, comprising the
13 remainder of the operating system 90. On the operating
14 system 90 is the application 600, comprising the application
15 software 80. Other stacked subdivisions of the data
16 processing system may be apparent to those skilled in the
17 art. In general, embodiments of the present invention may be
18 installed in one or more of the data processing layers
19 600-630. Each installation then protects the layer in which
20 it is hosted from corruption.

21 In the examples of the present invention herein before
22 described, all administrative actions producing a change in
23 state of the data processing system trigger an update of the
24 secret. However, in other embodiments of the present
25 invention, only a subset of such actions may trigger an
26 update of the secret. For example, only actions potentially
27 leading to infiltration of system level Trojan horses may
28 trigger an update of the secret.

29 Also, in the advantageous embodiments of the present
30 invention herein before described, the CEF 70 is implemented
31 by computer program code executing on the CPU 10 of the data
32 processing system. It will be appreciated that, in other
33 embodiments of the present invention, the CEF 70 may be

1 implemented at least partially by hardwired logic circuitry.
2 In particularly advantageous embodiments of the present
3 invention, the CEF 70 may be embedded in a trusted subsystem
4 of the data processing system.

5 Variations described for the present invention can be
6 realized in any combination desirable for each particular
7 application. Thus particular limitations, and/or embodiment
8 enhancements described herein, which may have particular
9 advantages to the particular application need not be used
10 for all applications. Also, not all limitations need be
11 implemented in methods, systems and/or apparatus including
12 one or more concepts of the present invention.

13 The present invention can be realized in hardware, software,
14 or a combination of hardware and software. A visualization
15 tool according to the present invention can be realized in a
16 centralized fashion in one computer system, or in a
17 distributed fashion where different elements are spread
18 across several interconnected computer systems. Any kind of
19 computer system - or other apparatus adapted for carrying
20 out the methods and/or functions described herein - is
21 suitable. A typical combination of hardware and software
22 could be a general purpose computer system with a computer
23 program that, when being loaded and executed, controls the
24 computer system such that it carries out the methods
25 described herein. The present invention can also be
26 embedded in a computer program product, which comprises all
27 the features enabling the implementation of the methods
28 described herein, and which - when loaded in a computer
29 system - is able to carry out these methods.

30 Computer program means or computer program in the present
31 context include any expression, in any language, code or
32 notation, of a set of instructions intended to cause a
33 system having an information processing capability to

1 perform a particular function either directly or after
2 conversion to another language, code or notation, and/or
3 reproduction in a different material form.

4 Thus the invention includes an article of manufacture which
5 comprises a computer usable medium having computer readable
6 program code means embodied therein for causing a function
7 described above. The computer readable program code means
8 in the article of manufacture comprises computer readable
9 program code means for causing a computer to effect the
10 steps of a method of this invention. Similarly, the present
11 invention may be implemented as a computer program product
12 comprising a computer usable medium having computer readable
13 program code means embodied therein for causing a function
14 described above. The computer readable program code means
15 in the computer program product comprising computer readable
16 program code means for causing a computer to effect one or
17 more functions of this invention. Furthermore, the present
18 invention may be implemented as a program storage device
19 readable by machine, tangibly embodying a program of
20 instructions executable by the machine to perform method
21 steps for causing one or more functions of this invention.

22 It is noted that the foregoing has outlined some of the more
23 pertinent objects and embodiments of the present invention.
24 This invention may be used for many applications. Thus,
25 although the description is made for particular arrangements
26 and methods, the intent and concept of the invention is
27 suitable and applicable to other arrangements and
28 applications. It will be clear to those skilled in the art
29 that modifications to the disclosed embodiments can be
30 effected without departing from the spirit and scope of the
31 invention. The described embodiments ought to be construed
32 to be merely illustrative of some of the more prominent
33 features and applications of the invention. Other
34 beneficial results can be realized by applying the disclosed

1 invention in a different manner or modifying the invention
2 in ways known to those familiar with the art.